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Routing in Wireless Mesh Networks

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About me


- **2005-:** PhD student at WIT (Kappel, IS)
 - Advisors: Dietrich (EE), Kastner (CS)
 - WP-Leader EU-Project Track & Trade
- **2001-2005:** Telecommunications Research Center Vienna (ftw)
- **Master:** Computer Science, Specialization: Computer Engineering
- **Research Interests:**
 - Wireless Computing and Telecommunications
 - Wireless Mesh / Ad-Hoc Networks
- **Thesis (working title):** "Routing in Wireless Mesh Networks"




Overview



- **My Thesis: "Routing in Wireless Mesh Networks"**
 - Definition
 - State-of-the-Art
 - Approach
 - Time Plan
- **Current Work**
 - Ns-2 Random Generator
 - Wireless Network Modelling

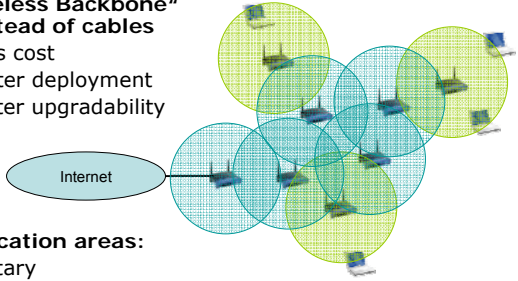


Definition / 1: Wireless Mesh Network





"Wireless Backbone" instead of cables

- Less cost
- Faster deployment
- Better upgradability


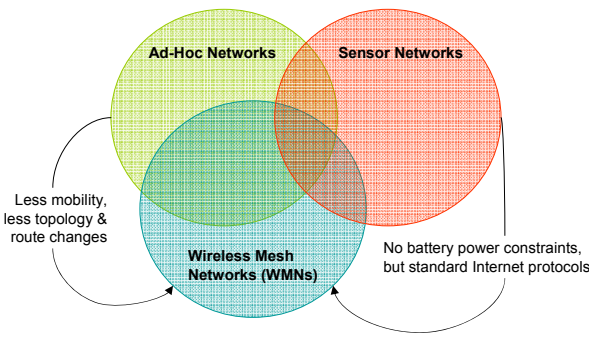


Application areas:

- Military
- Disaster areas / emergency response
- **Civilian: access network to the Internet**





Definition / 2: Field of Research





Less mobility, less topology & route changes

No battery power constraints, but standard Internet protocols

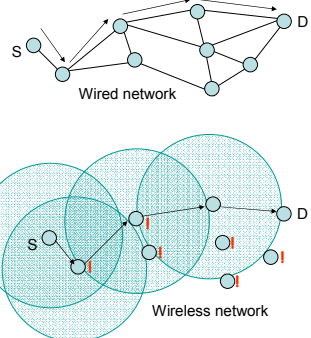


Definition / 3: The Problem




Routing:

- Hop-count is a bad metric:
 - Links vary over time
 - Far links are usually worse than short links



MAC Interaction:

- Hidden node problem
- Exposed terminal problem
- Flow-in-the-middle problem



State-of-the-Art / 1: Protocol types



- Reactive (on-demand)
 - Build route when flow starts
 - Destroy route when flow ends
 - Drawback: route setup delay
 - Good for frequent route changes
- Pro-active
 - Build route in advance
 - Keep routes up-to-date (communicate changes)
 - Drawback: communication overhead
 - Good when routes change seldom
- Hybrid



7

State-of-the-Art / 2: Protocols



"Classical":

- **OLSR** - Optimized Link State Routing protocol pro-active, *RFC 3626* (Clausen, Jaquet, 2003), used in firmwares
- **AODV** - Ad-hoc On-Demand Distance Vector reactive, *RFC 3561* (Perkins, Royer, Das, 2003)
- **DSDV** - Dynamic Destination-Sequenced Distance Vector pro-active, one of the oldest (Perkins, 1994)
- **DSR** - Dynamic Source Routing reactive (Johnson, 1994)

Nature inspired:

- Eg. Ant Hoc Net hybrid (DiCaro et al, 2004)



8

Ant-based Algorithms / 1



- Inspired by nature: behavior of ants
- Single ants are quite stupid, but the whole system exhibits "intelligent" behavior
- Ant Colony Routing (ACR) – distributed version of Ant-based Algorithm, eg.: **AntNet** by Di Caro and Dorigo, 1998
- **AntHocNet** for MANETs by Di Caro, Ducatelle, Gambardella, 2004:
 - AntNet concept + Extensions
 - Hybrid routing approach: reactive/pro-active



9

Ant-based Algorithms / 2



1. Whenever an ant moves, it lays a pheromone trail
 2. To find its way, an ant:
 - Follows existing trails if there are any. Probability for choosing a trail is proportional to amount of pheromone on the trail.
 - Walks randomly if there are no trails.
 3. Pheromone evaporates over time -> unused trails vanish.
- **Trail following (state transition rule)** determines how the ant chooses its way depending on link cost and amount of pheromone found on the trails
 - **Trail laying (pheromone update rule)** determines how the pheromone is updated
 - **Evaporation (evaporation rule)** determines how fast pheromone evaporates



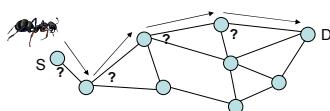
10

Ant-based Algorithms / 3



- Forward ants: regularly created, choose next link based on transition rule:

$$P = \text{Trans}(\text{pheromone}, \text{link cost})$$

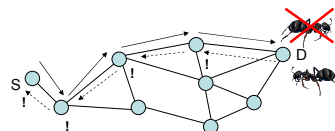


11

Ant-based Algorithms / 4



- Backward ants: created when forward ant reaches destination node, travels back to from where it came, updates pheromone amounts and measured link costs at all nodes on way back



12

My Approach

- Simulation with ns-2 simulator
- An implementation of AntNet has just become available

Find new algorithm:

- Based on ant concept
- Probably hybrid between classical and ant-based
- Probably hybrid proactive/reactive routing approach – WMNs are less mobile than MANETs
- Use better cross-layer info to determine link costs
- Reduce number of ants by observing TCP traffic in the network (possible?)
- Idea: use "colored pheromones" to signify QoS classes -> different traffic takes different routes in the network

Time Plan

- 25.7.:** Student poster at INDIN
- 1.8.:** Book chapter of ns-2 RNG paper
- 5.8.:** Define and implement algorithm in ns2; paper for BIONETICS 2007, Budapest
- 21.9.:** Evaluation of wireless models in ns-2; SIMUTools 2008, Marseille
- 7.-9.11.:** FET (Toulouse) – short paper (wireless modelling) under review
- 12/2007:** performance evaluation; publish 2 papers about algorithm
- 05/2008:** Finish Thesis

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Current Work

- Ns-2 Random Generator
- Wireless Network Modelling

Current Work: Ns-2 Random Generator

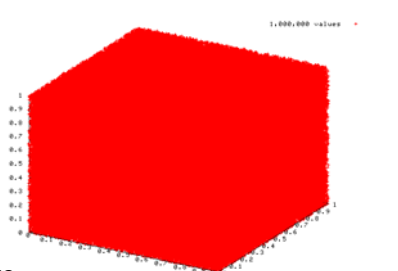
- Martina Umlauf, Peter Reichl, *Experiences with the ns-2 Network Simulator - Explicitly Setting Seeds Considered Harmful*
- Presented at Wireless Telecommunications Symposium WTS '07, 26.-28. April, Pomona, CA, USA
- Selected to be published as book chapter in *Wireless Technology: Applications, Management, and Security*, Springer

The ns-2 RNG

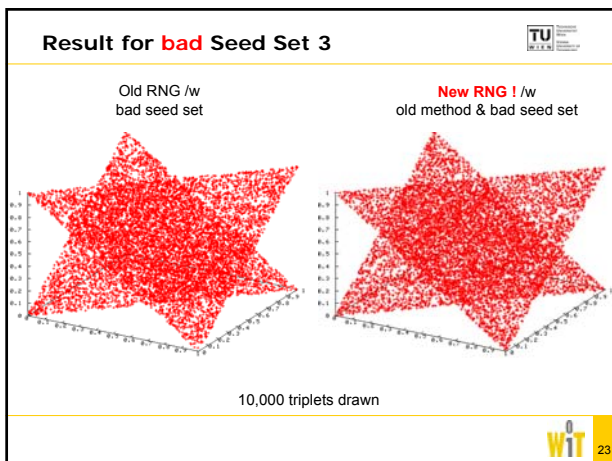
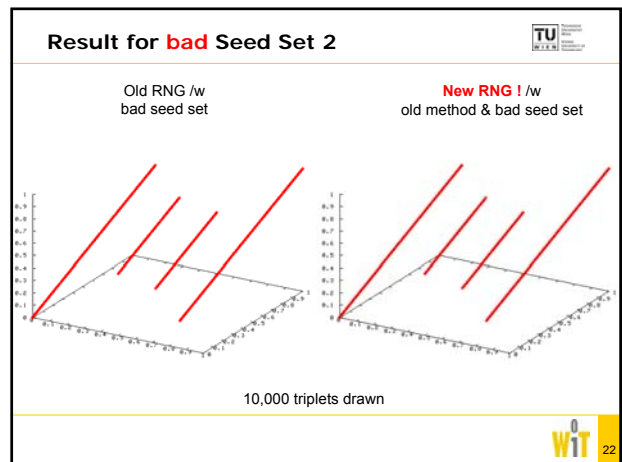
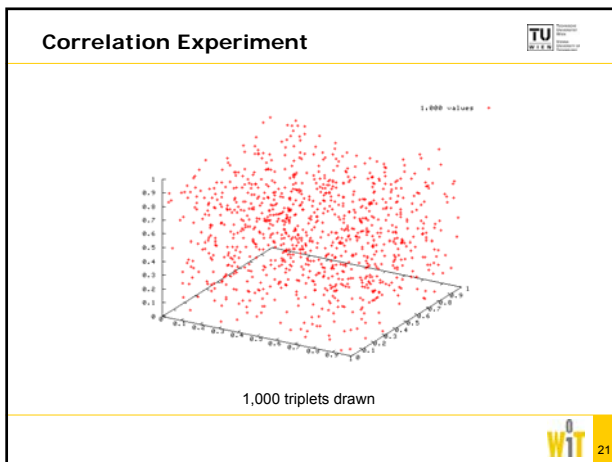
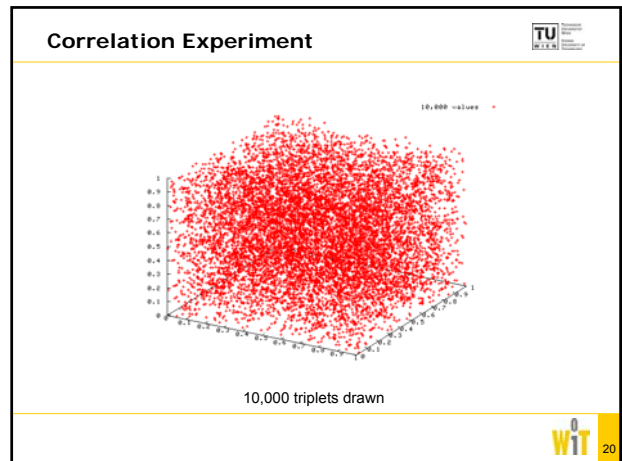
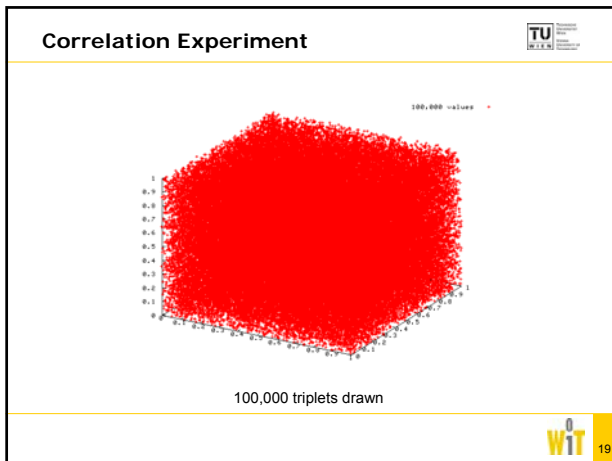
- ns-2 versions $\leq 2.1b8$: **old RNG**
Minimal Standard multiplicative Linear Congruential Generator [Park and Miller, 1988]
Period $p=2^{31}-2$
API: `$rng seed $s;`
Sensitive to seeds [Entacher, Hechenleithner, 2002]
- ns-2 versions $\geq 2.1b9$ until today: **new RNG**
MRG32k3a [L'Ecuyer, 1999]
Period $p = 3.1 \times 10^{57}$
API: `$rng next-substream;`
Promises to fix seed-sensitivity -> **true?**

BUT: bad documentation -> people use old API!

Correlation Experiment



- Set up 3 RNGs**
- Draw triplets of values**
`$u($i)` from `$rng($i)` 1,000,000 triplets drawn
- Interpret as vector**
`<$u(1), $u(2), $u(3)>`



Wired Topology Example / 1

Exponential On/Off-Traffic: $\lambda = 8000\text{bits} / 41\text{ms} = 0.195\text{Mbps}$
 On: $0.08\mu\text{s} / 1\text{pkt} / 1000\text{bytes}$
 Off: 41ms

$\Sigma\lambda = 0.976\text{Mbps}$
 $\rho = \Sigma\lambda / \text{BW} = 0.976 \dots \text{utilization factor}$

Expected mean queue length:
 $\bar{q} = \frac{\rho}{1-\rho} - \frac{\rho^2}{2(1-\rho)} = 20.488\text{pkts}$

New RNG	New method	20.2996
New RNG	Set 2 - bad	29.4527
Old RNG	Set 1 - good	19.4398
Old RNG	Set 2 - bad	24.2785

Simulation time 7200s
 Sampling time 10ms

Bad seed sets: higher average queue lengths!

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Wireless Example / 1 TU
TU Braunschweig

Simulation time 600s

Run lengths vs. burst lengths:

Bad seed sets: higher percentage of short run lengths!

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Overview TU
TU Braunschweig

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Wireless Network Modelling

WIT 26

Current Work: Wireless Network Modelling TU
TU Braunschweig

- Martina Umlauf, *A Better Way to Model Wireless Networks*, short paper, submitted to FET '07
- Martina Umlauf, *Some Thoughts on Wireless Network Modelling*, Student Poster, accepted at INDIN '07, to be presented 2007-07-25

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Traditional Reachability Graph G TU
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- Edges $L_{1,3}$ and $L_{4,8}$ seem independent
- Edges encourage thinking of links as "tunnels" between nodes

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Vs. Real World Situation TU
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Theoretical transmission range

Carrier sense (CCA) range (idealized)

Example of real-world signal propagation as measured in (Kotz et al, 2003)

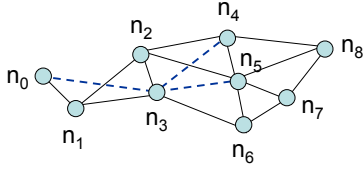
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Some Observations TU
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- No "tunnel" - so what is a "link"?
 - Unicasts "do not exist" - can lead to: **hidden node problem, exposed node problem**
 - (Padhye et al, 2005):
 $\exists L_{i,j}$ in G iff: $n_i \rightarrow n_j \cap n_i \rightarrow n_j$
 with packet loss rate below certain threshold ($ETX \leq 3$)
 ETX (DeCouto et al, 2003)
- Carrier sensing (CCA)
 - A node can "hear" the transmission but SINR is too bad to successfully decode the packet
 - CCA function: channel is busy! Do not send (**exposed node problem**)
- "Layer 2 interference" vs. "Interference" on layer 1

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Proposal: new graph G^* with CCA-edges

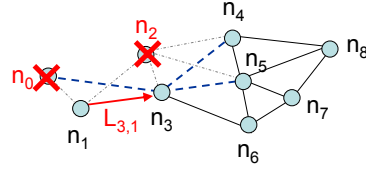


Only one paper found which informally uses such a graph (Das et al, 2006).



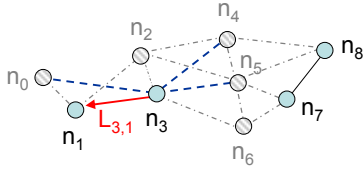
31

Transmission of Data Packet $n_1 \rightarrow n_3$



32

Transmission of ACK $n_3 \rightarrow n_1$



33

References / Wireless Network Modelling



- Das et al, 2006: **Characterizing Multi-Way Interference In Wireless Mesh Networks**. In *Proc. 1st Intl. Workshop on Wireless network testbeds, experimental evaluation & characterization (WINTECH '06)*, Los Angeles, CA, USA, Sept. 2006.
- DeCouto et al, **A High-Throughput Path Metric for Multi-Hop Wireless Routing**. In *Proc. ACM MobiCom*, San Diego, CA, USA, Sept. 2003.
- Kotz et al, 2003: **The mistaken axioms of wireless network research**. Technical Report TR2003-467, Dartmouth College of Computer Science, 2003.
- Padhye et al, 2005: **Estimation of Link Interference in Static Multi-Hop Wireless Networks**. In *Proc. ACM/USENIX Intl. Measurement Conference*, Berkeley, CA, USA, Oct. 2005.



34

Thank You!



35